

# FireWire Memory Dump of a Windows XP Computer: A Forensic Approach

Antonio Martin  
[tmemail@gmail.com](mailto:tmemail@gmail.com)  
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## Introduction

In a forensic investigation, while collecting evidence, Anzaldua, Godwin, Volonino state the best practice is to unplug a computer or remove a laptop's battery, so as to preserve the exact contents of a disk without introducing artifacts, changes to the system, for later investigations.<sup>[1]</sup> It is critical the original contents of a system, as found, are not altered. Using the operating system shut down alters log and temp file states; furthermore, a shutdown may trigger a logic bomb and a possible data wipe. This approach has a problem; it does not preserve active memory. A portion of the active memory can be found in the remnants of the operating system swap file, but this is an incomplete picture. There exists a desire to find a means to collect a forensic image of memory without compromising an investigation's integrity.

## Network Based, Active Data Collection

As the field of digital forensics progresses, new means of evidence collection emerge. The current “best practice” allows for *In situ*, on sight, live data collection from running systems; a snapshot capture of all the information in memory. While desirable, it results in a series of problems that have not been addressed. Is this a case of law and forensic practitioners being technologically a few steps behind?

One example of the current means of network based, live data collection can be found in Guidance Software’s Encase® Field Intelligence Model (EFIM); other software packages behave similarly. EFIM allows a forensics investigator to connect to a target machine by Ethernet, push a small program referred to as a “servlet” and capture the system’s live memory and hard drive. EnCase

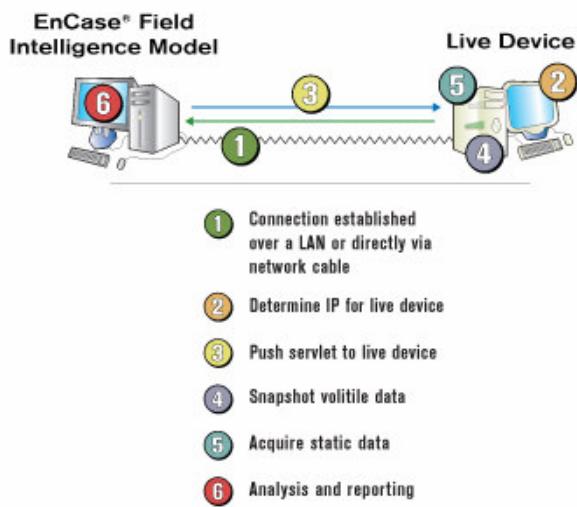


Figure 1. From [www.guidancesoftware.com](http://www.guidancesoftware.com)

FIM is capable of targeting “Windows 95/ 98/ NT/ 2000/ XP/ 2003 Server, Linux Kernel 2.4 and above, Solaris 8/9 both 32 & 64 bit, AIX, OSX.”<sup>[2]</sup> At first glance, as long as the artifacts that are introduced by the “servlet” are well understood, the evidence collected should be trustable and admissible. The tool allows investigators to image running systems and servers while not disrupting a company’s operations. Furthermore, it can be used to find and target machines on a network (wirelessly?), without the owner’s knowledge.

EnCase’s approach presents at least the following issues:

- A “servlet” installation introduces artifacts on the target system. Claims that the impacts are “well documented and known” are interesting; it is impossible to test and document all possible hardware and software configurations and how their interactions will affect the impacts of installation and operation.
- A “servlet” is susceptible to attack since malicious software running on the target can identify and stop it or trigger a logic bomb. This becomes more of an issue since the full Encase product suite is available from torrents and warez sites allowing criminals to dissect and build defenses. While the “servlets” can be updated and altered, it will require investigators to constantly update and document changes to their investigation test stands and run the risk of losing all data if a logic bomb is triggered. Simple file

modification of the “servlet” might not work to prevent detections as data stream and behavior analysis can quickly identify potential alterations/updates.

- It is possible the “servlet” could be maliciously submitted to malware and virus protection houses. The code would be inspected and signature detection profiles pushed out to millions of computers world wide. Thus a target with active, running virus protection might automatically stop a live forensics investigation.
- By far, the most critical issue: Any targeted system is an unknown quantity with little insight as to what is or is not running. If a system is infected, it is impossible to trust any information gathered through the operating system because stealth root kits are difficult, if not impossible to detect. Rootkits have advanced considerably and data from an unknown, possibly infected system cannot be trusted. <sup>[3, 4]</sup> A kernel mode rootkit called “Shadow Walker,” available since summer of 05, is such a program; it hides itself and other processes within memory by subverting the windows virtual memory management. Any process from within the CPU attempting to view memory will be fooled.

*“Unfortunately, because all live response tools of which we are aware run directly on a potentially compromised system, they rely on the underlying operating system, .... Even tools which attempt to determine the integrity of the operating system may be fooled if the attacker has perfect knowledge of the tool and control of the system before it is installed. ... it is impossible to even know if the live response tool itself has been run in an unmodified way. This means that, even if the tool itself has been verified, the executing instance of that tool may be untrustworthy.”* <sup>[Butler, Sparks 4]</sup>

It is possible that a rootkit could be discovered on analysis of the hard drive copy, if the rootkit allowed the drive section to be read. The rootkit might not exist in the drive, but could hide its image in the BIOS or a video card ROM.<sup>[3]</sup>

The process and end product of gathering forensic evidence is supposed to be of the utmost quality and integrity, but this approach to live evidence collection is flawed and approaches negligence if it makes such claims. Ways to subvert and bypass such methods are already used and well documented. While this can be a useful tool for information gathering, this form of forensic evidence collection should not be admissible in a court of law where integrity of data is of the utmost concern because it cannot be guaranteed.

## Firewire 1394 or or

Firewire® is a bus technology designed for point to point connections between devices. It was developed by Apple in 1986 and was standardized to the IEEE 1394 specification in 1995.<sup>[5]</sup> Firewire, like many other devices in today's computer architecture, utilizes Direct Memory Access (DMA) to improve data transfers.

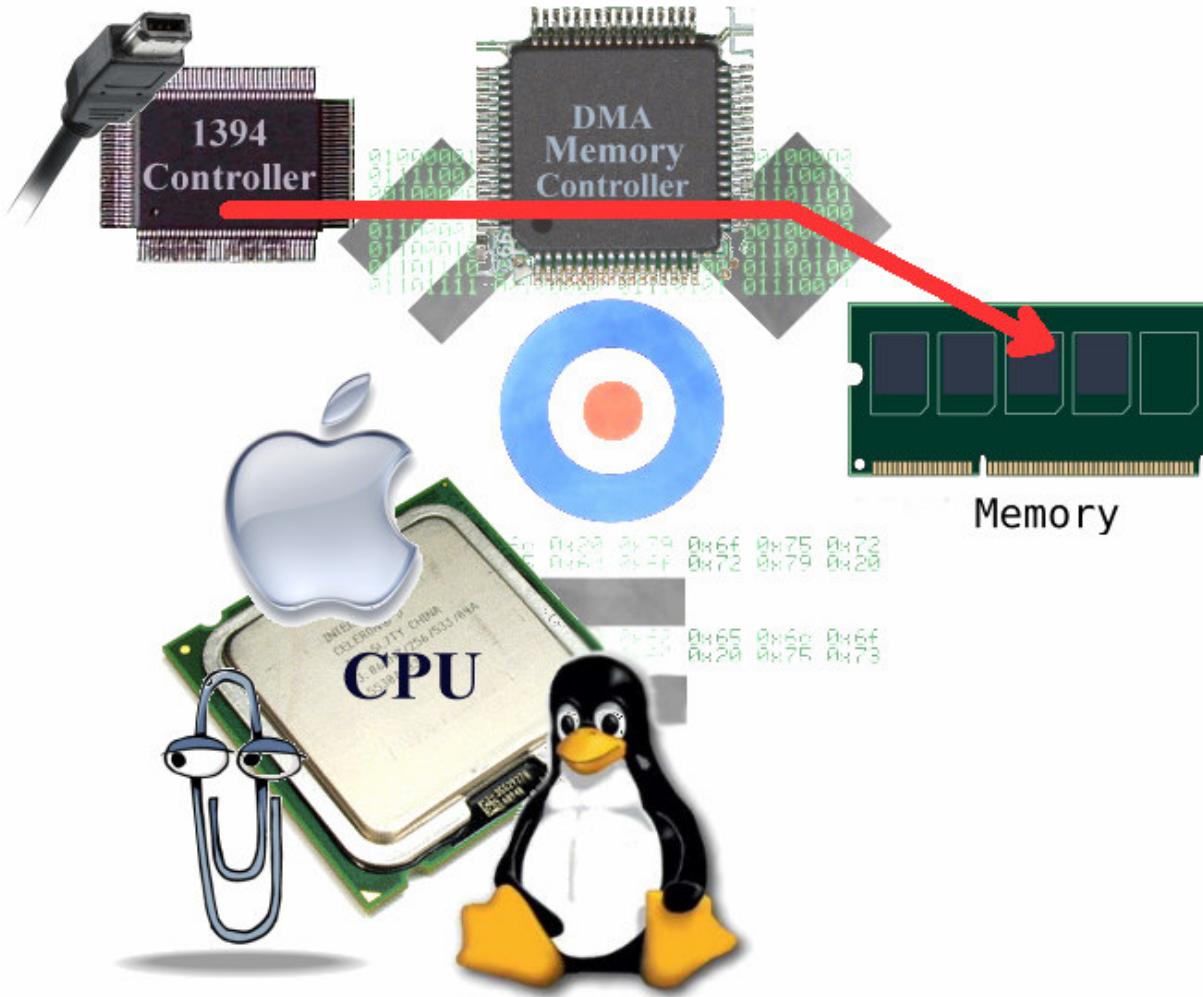


Figure 2: Firewire, using DMA, bypasses the CPU and running operating system.

A firewire device can read (and write) to a computer's main memory by accessing a system's DMA controller, while the operating system, be it Windows, Mac OS, Linux, a Multiple Independent Levels of Security kernel, etc., is oblivious to the event. By pulling a copy of memory through firewire, the target CPU and operating system are bypassed as are any infections, triggers or traps. This is not a bug but exactly how DMA and PCI devices, like Firewire, were designed to operate.

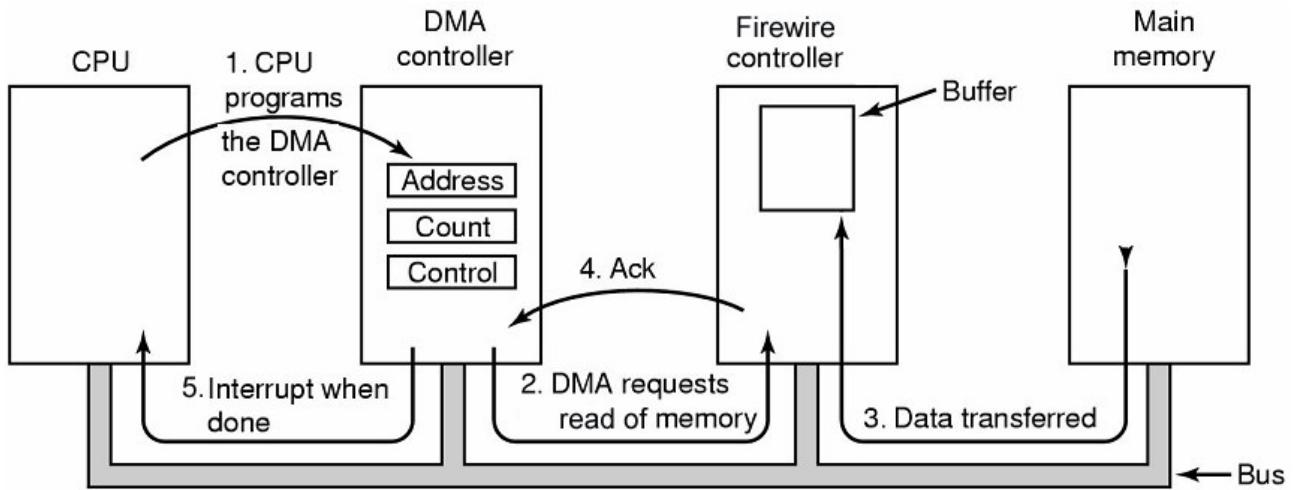


Figure 3: Modified from IO class notes, Dr Szabolcs Mikulas, School of Computer Science and Information Systems, Birkbeck College, Chapter 5, Input / Output [6]

DMA allows memory transfers between devices and processes to take place while a computer's CPU performs other tasks. Figure 3:

1. The CPU/operating system programs the DMA controller to instruct a Firewire device to read a portion of memory; the CPU/operating system is now free to work on other tasks.
2. The DMA controller sends a message to the Firewire controller, informing it of the read request and the location and length in memory.
3. The Firewire device negotiate control of the PCI bus and reads the memory location specified and once completed,
4. Informs the DMA controller.
5. Finally, the DMA controller triggers an interrupt, informing the CPU the read has completed.

It should be noted that devices are not limited to only reading/writing to the memory address specified by the operating system. Firewire and other DMA bus master devices act independently of the CPU; the CPU need not initiate the transaction. A firewire device can program the DMA controller and set up its own reads and writes, as per the PCI and IEEE 1394 specifications.

## “Hit by a Bus: Physical Access Attacks with Firewire”

At PacSec in November of 2005, Maximillian Dornsief presented a paper, “*Owned by an iPod*” where he demonstrated how a firewire device, utilizing DMA, can read/write active memory within a Mac, BSD or Linux machine.<sup>[7]</sup> At Ruxcon 2006, Adam Boileau (aka Metlstorm) from Security-Assessment.com presented “*Hit By A Bus: Physical Access Attacks with Firewire*” where he extended prior works, enabling the targeting of a Windows XP machine. Utilizing a Linux box

with firewire support, and a set of tools for enabling the interface, he revealed several hacks using live memory reading and writing targeted against Windows XP.

- Reading over firewire the entirety of the target's memory and saving it to disk without altering the target's state.
  - Locating and over writing a memory address containing the graphical identification and authentication library (msgina.dll) allowing the password on a locked Windows XP machine to be bypassed.
  - Pushing to the target computer and starting a process without the process existing on the target hard drive.
  - Recovering the last sixteen bytes from the keyboard buffer accepted by the BIOS prior to booting the primary operating system, useful in finding BIOS and disk encryption passwords.

Along with the presentation was released a set of Linux tools for firewire memory reading and writing.<sup>[8]</sup>

# What is possible?

Utilizing MetlStorm's toolset, it is possible to configure a Linux system (in this case Ubuntu 6.10) and target an IBM Trusted Computing Module enabled Thinkpad.

000000600	FA 33 C0 8E C0 8E D8 8E D0 BC 00 7C 8B F4 FB FC	ü3A À Ø Ð¾.   òüü
000000610	BF 00 06 B9 00 01 F3 A5 B8 DF 06 50 C3 00 0F 00	ë..¹..ó¥,ß.PÃ...
000000620	01 0A 45 72 72 6F 72 20 6C 6F 61 64 69 6E 67 20	.Error loading
000000630	6F 70 65 72 61 74 69 6E 67 20 73 79 73 74 65 6D	operating system
000000640	0A 0D 00 0A 49 6E 76 61 6C 69 64 20 70 61 72 74	....Invalid part
000000650	69 74 69 6F 6E 20 74 61 62 6C 65 0A 0D 00 50 72	ition table...Pr
000000660	65 73 73 20 6B 65 79 20 74 6F 20 72 65 62 6F 6F	ess key to reboo
000000670	74 20 0A 0D 00 80 7C 04 0C 74 1A 80 7C 04 0E 74	t ...  ...t...  ...t
000000680	14 81 7C 0A FA 00 73 0D 8B 4C 02 8B 14 B2 80 B8	.  .ú.s. L ..² ,
000000690	01 02 CD 13 C3 56 8B C3 87 DE BE 00 06 C7 04 10	. ÄV Ä Þ¾..ç..
0000006A0	00 C7 44 02 01 00 89 44 04 8C 5C 06 8B 47 08 89	.QD... D. N. G.
0000006B0	44 08 8B 47 0A 89 44 0A C7 44 0C 00 00 C7 44 0E	D. G. D.CD...CD.

*Figure 4: Memory dump containing a system's BIOS.*

The firewire dump produces a large (the size of available memory) and difficult to decipher binary image. Figure 4 shows a portion of the capture displaying the target's BIOS in memory. This file is difficult to understand without in-depth knowledge of the target operating system's memory map; thankfully there are several tools available to assist. Andreas Schuster created Process and Thread Finder (PTFinder), a script capable of parsing firewire memory dumps.<sup>[9]</sup> With this script and a little work, the following information can be gathered:

An investigator can quickly know what programs are/were running and can view the various program/thread memories for information about activity, connections, username password combinations, etc. Below is an example of a PTFinder list from the IBM memory dump. It dumps all threads (and processes), their thread IDs, associated Process IDs, times created, exited, offset/location within the memory dump (so you can go to the thread's location and view the information), the PDB (processes virtual address value) and Remarks (usually the process name or system's status).

No.	Type	PID	TID	Time created	Time exited	Offset	PDB	Remarks
1	Thrd	0	0			0x00559320		
2	Proc	0				0x00559580	0x00039000	idle
3	Thrd	4	3284		2006-08-03 09:12:35	0x02a66da8		
4	Thrd	4	3496		2006-08-03 16:43:45	0x02a80da8		
5	Thrd	4	4048		2006-08-03 16:43:45	0x02a81da8		
6	Thrd	4	3088		2006-08-03 16:43:42	0x02a82da8		
7	Thrd	4	3412		2006-08-03 16:43:42	0x02a83020		
8	Thrd	4	3572		2006-08-03 16:43:42	0x02a833c8		
9	Thrd	4	2660		2006-08-03 16:43:42	0x02a83640		
10	Thrd	4	4036		2006-08-03 16:43:42	0x02a838b8		
...								
649	Proc	2368			2006-07-31 16:18:15	0x0368ba98	0x16ef8000	ibmmessages.exe
695	Proc	4				0x037c87c0	0x00039000	System

Running a grep on the list or using a flag for PTFinder can produce a list of just the processes running that were found in the memory dump. Note the first two, red highlighted lines (No. 55 and 247), a Back Orifice 2k configuration tool (bo2kcfg.exe) and the BO2K GUI (bo2kgui.exe) used to control BO2K infected machines. In some cases, the residual memory from prior processes and threads is still available; the third highlighted line (processes SynTPLre.exe, No. 292) shows a recently exited process whose memory is still available.

No.	Type	PID	TID	Time created	Time exited	Offset	PDB	Remarks
2	Proc	0				0x00559580	0x00039000	idle
55	Proc	2060		2006-08-03 16:32:39		0x02d8d020	0x24445000	bo2kcfg.exe
148	Proc	3388		2006-08-03 16:51:05		0x02e1c4c0	0x28659000	msmsgs.exe
247	Proc	624		2006-08-03 16:32:41		0x02ee0020	0x17a4f000	bo2kgui.exe
254	Proc	2392		2006-07-31 16:18:17		0x02eedda0	0x16c57000	tfsctrl.exe
292	Proc	2156		2006-07-31 16:18:11	2006-07-31 16:18:12	0x02f2fb00	0x14baf000	SynTPLpr.exe
293	Proc	2056		2006-07-31 16:18:09		0x02f31800	0x1499e000	Tpscrex.exe
295	Proc	1972		2006-07-31 16:18:08		0x02f33d00	0x14597000	TPONSCR.exe
297	Proc	2480		2006-07-31 16:18:20		0x02f38638	0x1755d000	certtool.exe
299	Proc	2528		2006-07-31 16:18:22		0x02f3bd00	0x1789e000	pwmgr.exe
308	Proc	724		2006-07-31 16:18:06		0x02f4c608	0x1408c000	QCWLICON.EXE
321	Proc	968		2006-07-31 16:18:04		0x02f59620	0x13d54000	TPHKGMR.exe
329	Proc	1804		2006-07-31 16:17:53		0x02f619e0	0x1323a000	explorer.exe
344	Proc	2588		2006-07-31 16:18:22		0x02f77488	0x17ead000	ccApp.exe
349	Proc	224		2006-07-31 16:15:10		0x02f7d630	0x0e557000	a1g.exe
362	Proc	1904		2006-07-31 16:15:08		0x02f8ada0	0x0db41000	SymWSC.exe
371	Proc	1844		2006-07-31 16:15:07		0x02f96da0	0x0d939000	TpkmpSvc.exe
389	Proc	1708		2006-07-31 16:15:06		0x02faf020	0x0d129000	RegSrvc.exe
390	Proc	1744		2006-07-31 16:15:07		0x02faf800	0x0d22e000	UMGR32.EXE
393	Proc	1688		2006-07-31 16:15:06		0x02fb4800	0x0d10b000	TssCore.exe
406	Proc	1636		2006-07-31 16:15:06		0x02fc8da0	0x0d122000	QCONSVC.EXE
419	Proc	896		2006-07-31 16:18:02		0x02fd8948	0x14116000	SynTPEnh.exe
427	Proc	1592		2006-07-31 16:15:06		0x02fe2800	0x0cc1b000	uvmserv.exe
430	Proc	1544		2006-07-31 16:15:06		0x02fe3bc0	0x0ce16000	ati2evxx.exe
439	Proc	1440		2006-07-31 16:15:05		0x02feebc0	0x0cc08000	spoolsv.exe
463	Proc	1308		2006-07-31 16:15:05		0x03003b88	0x0c1fe000	ccEvtMgr.exe
479	Proc	1764		2006-07-31 16:18:05		0x0300cd40	0x13b76000	TP98TRAY.EXE
485	Proc	1128		2006-07-31 16:15:04		0x030111da0	0x0c0f6000	svchost.exe
492	Proc	976		2006-07-31 16:15:03		0x03018800	0x0bbd2000	svchost.exe

494 Proc 492	2006-07-31 16:18:02	0x0301ada0 0x13dd0000	SynTPLpr.exe
519 Proc 1084	2006-07-31 16:15:04	0x032098b0 0x0c0f0000	svchost.exe
523 Proc 876	2006-07-31 16:15:01	0x0320c648 0xb79d000	svchost.exe
535 Proc 1032	2006-07-31 16:15:04	0x0321f020 0x0bfe5000	S24EvMon.exe
557 Proc 840	2006-07-31 16:14:58	0x03380da0 0x0b696000	ibmpmsvc.exe
563 Proc 656	2006-07-31 16:14:57	0x033b0bc0 0x0aae0000	lsass.exe
579 Proc 644	2006-07-31 16:14:57	0x033f29c0 0xa9d6000	services.exe
582 Proc 600	2006-07-31 16:14:55	0x03400da0 0xa9b4000	winlogon.exe
608 Proc 556	2006-07-31 16:14:48	0x034f9da0 0x09dae000	cssrss.exe
621 Proc 940	2006-07-31 16:15:02	0x03559600 0xb9cb000	svchost.exe
637 Proc 500	2006-07-31 16:14:40	0x035b23e0 0x08943000	smse.exe
644 Proc 2188	2006-07-31 16:18:12	0x03680da0 0x15bfc000	AGRSMMSG.exe
646 Proc 2348	2006-07-31 16:18:15	0x03688aa8 0x164c3000	tgcmd.exe
647 Proc 2176	2006-07-31 16:18:12	0x036894f8 0x14976000	EzEjMnAp.Exe
649 Proc 2368	2006-07-31 16:18:15	0x0368ba98 0x16ef8000	ibmmessages.exe
695 Proc 4		0x037c87c0 0x00039000	System

Six hundred and ninety five threads were identified and notated in the memory dump by PTFinder. A useful enhancement to the PTFinder tool would be the ability to save the memory sections for the individual processes and associated threads, each stored in their own file grouped by directory. This would allow quicker and easier examination and categorization.

02D4C720	4C 4D 45 4D 28 00 00 00	FC 67 17 00 78 36 16 00	LMEM(...üg..x6..
02D4C730	47 00 08 00 B1 01 0A 00	A8 19 18 00 6F 72 69 7A	G...±....oriz
02D4C740	61 74 69 6F 6E 3A 20 50	61 73 73 70 6F 72 74 31	ation: Passport1
02D4C750	2E 34 20 4F 72 67 56 65	72 62 3D 47 45 54 2C 4F	.4 OrgVerb=GET,O
02D4C760	72 67 55 52 4C 3D 68 74	74 70 25 33 41 25 32 46	rgURL=http%3A%2F
02D4C770	25 32 46 6D 65 73 73 65	6E 67 65 72 25 32 45 6D	%2Fmessenger%2Em
02D4C780	73 6E 25 32 45 63 6F 6D	2C 73 69 67 6E 2D 69 6E	sn%2Ecom,sign-in
02D4C790	3D 6E 6F 61 63 63 6F 75	6E 74 31 32 33 34 25 34	=noaccount1234%4
02D4C7A0	30 68 6F 74 6D 61 69 6C	2E 63 6F 6D 2C 70 77 64	Ohotmail.com.pwd
02D4C7B0	3D 70 61 73 73 77 72 64	2C 6C 63 3D 33 33 33 33	=passwd1lc=3333
02D4C7C0	2C 69 64 3D 33 33 33 2C	74 77 3D 33 33 2C 72 75	,id=333,tw=33,ru
02D4C7D0	3D 68 74 74 70 25 33 41	25 32 46 25 32 46 6D 65	=http%3A%2F%2Fme
02D4C7E0	73 73 65 6E 67 65 72 25	32 45 6D 73 6E 25 32 45	ssenger%2Emen%2E
02D4C7F0	63 6F 6D 2C 63 74 3D 31	31 37 35 35 31 39 39 38	com,ct=117551998
02D4C800	37 2C 6B 70 70 3D 31 2C	6B 76 3D 39 2C 76 65 72	7,kpp=1,kv=9,ver
02D4C810	3D 32 2E 31 2E 36 30 30	30 2E 31 2C 72 6E 3D 61	=2.1.6000.1,rn=a

Figure 5: Memory section for Windows Messenger with username and password to the account.

Opening the saved memory image file in a hex editor (WinHex used) allow an examiner to find the memory sections pointed to by the PTFinder dump. Referencing the Windows Messenger (msmsg.exe) process id 3388, it is possible to find all associated threads in the dump and go to those memory offsets. Figure 5 shows an example thread 2248's memory section that contains the Windows Messenger's sign-in id and password.

17B7AC00	00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00	.....
17B7AC10	0C 00 00 00 4C 4F 47 4F 4E 5F 4F 42 4A 45 43 54	....LOGON_OBJECT
17B7AC20	30 1C 3E 02 05 00 00 00 00 00 00 05 00 00 00 00	0.>.....
17B7AC30	44 61 76 65 00 00 46 00 00 00 00 00 0E 00 00 00	Dave..F.....
17B7AC40	53 59 53 54 45 4D 5F 55 4E 4C 4F 43 4B 00 46 00	SYSTEM_UNLOCK.F.
17B7AC50	0C 01 00 00 5C 1C 3E 02 0C 01 00 00 03 00 00 00 00	....\>.....
17B7AC60	50 61 73 73 77 6F 72 64 31 31 00 00 00 00 00 00	Password11.....
17B7AC70	00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00	.....

Figure 6: Memory dump from an IBM Thinkpad with TMP, user ID and password readable.

Looking into kernel memory can reveal interesting information. Figure 6 displays a portion of memory taken from an IBM (now Lenovo) Thinkpad with a Trusted Computing Module and associated software. A TMP laptop is marketed offering a higher level of security over a standard laptop by leveraging TMP protections. This structure was found in high memory and contains the username, Dave, and password, Password11, of the person currently logged into the targeted machine. While a TPM system is supposed to operate at a higher level of security, the reality is much different.

The set of firewire tools for creating the memory images from MetlStorm has been added to the FCCU GNU/Linux Forensic Boot CD and can be found at <http://www.lnx4n6.be/>.<sup>[10]</sup>

Konuku's Volatools also offers a set of tools for analyzing memory images but it appears to not be designed for firewire dumps.<sup>[11]</sup> It failed on many of the attempts to parse most information from the file, like processes and threads but was able to find the current computer time.

```
E:\Python25> python volatools ident -f memoryimage.bin
Image Name: memoryimage.bin
Image Type: XP SP2
VM Type: nopae
DTB: 0x39000
Datetime: Thu Aug 03 09:51:16 2006
```

It also has support to find open sockets and network connection addresses (also failed). A potentially useful tool in analyzing Windows memory dumps.

## Can it be defeated?

At the February 2007 BlackHat convention in DC, Joanna Rutkowska demonstrated how to defeat DMA based memory gathering by utilizing a low level program, in the CPU, to rewrite a computer's North Bridge memory lookup table's pointers.<sup>[12]</sup> The north bridge hosts its own memory lookup table (IO Memory Management Unit), a mapping of the addresses and layout in main memory. By rewriting the lookup table, three possible means of preventing firewire/DMA

based memory reads can be realized. The first redirects the IO back at the system bus; this causes the computer to freeze, probably because the address range was not valid. The second redirects the IO so that all data returned is 0xff. Lastly, a more stealth means where certain memory pages can be hidden by removing their pointers and addressing them to other locations.

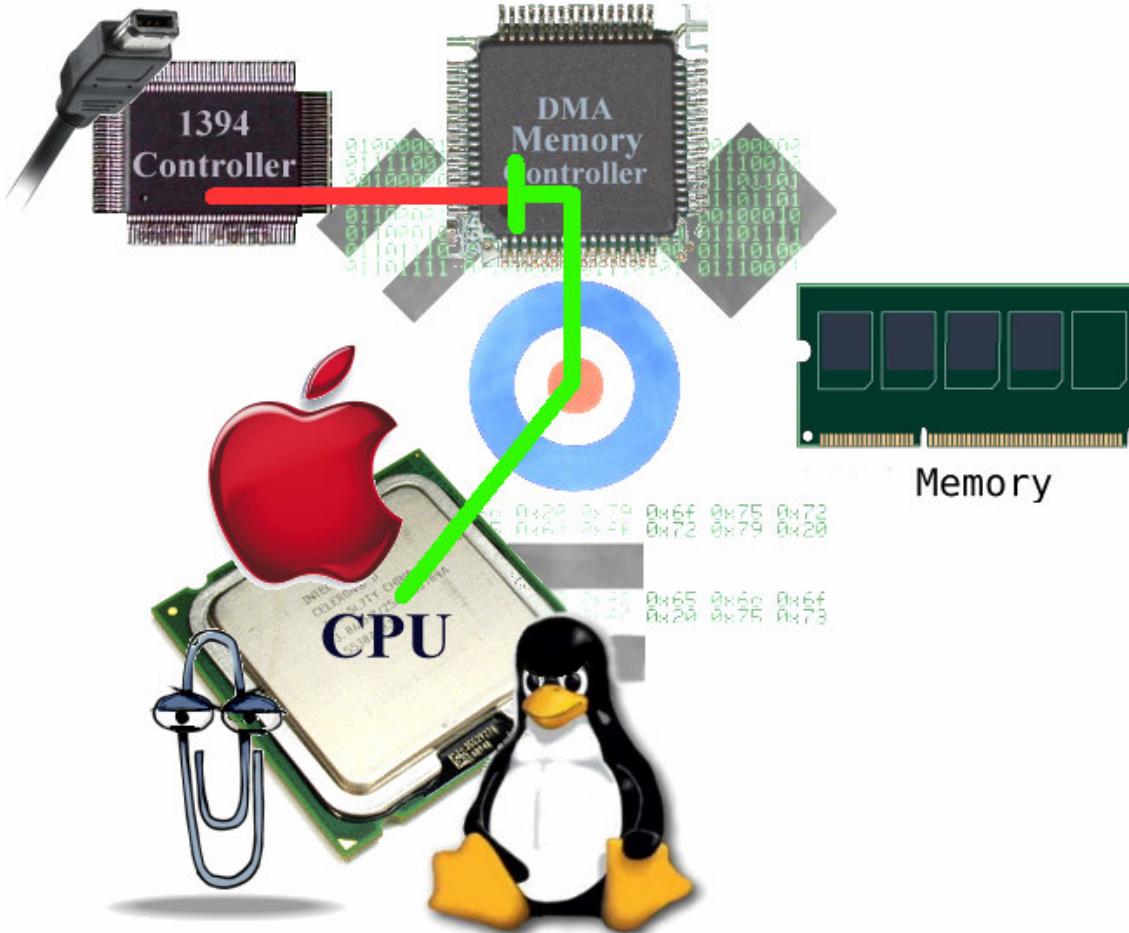


Figure 7: The operating systems get mean, remapping the DMA address table, preventing DMA.

# Imaging a Live Drive by Firewire?

Apple's computers offer a "Target Disk Mode," providing the ability for one Mac computer to boot off of another's drive by a firewire connection. This leads to the possibility of not only collecting active memory from a running system but also the contents of the hard drive.

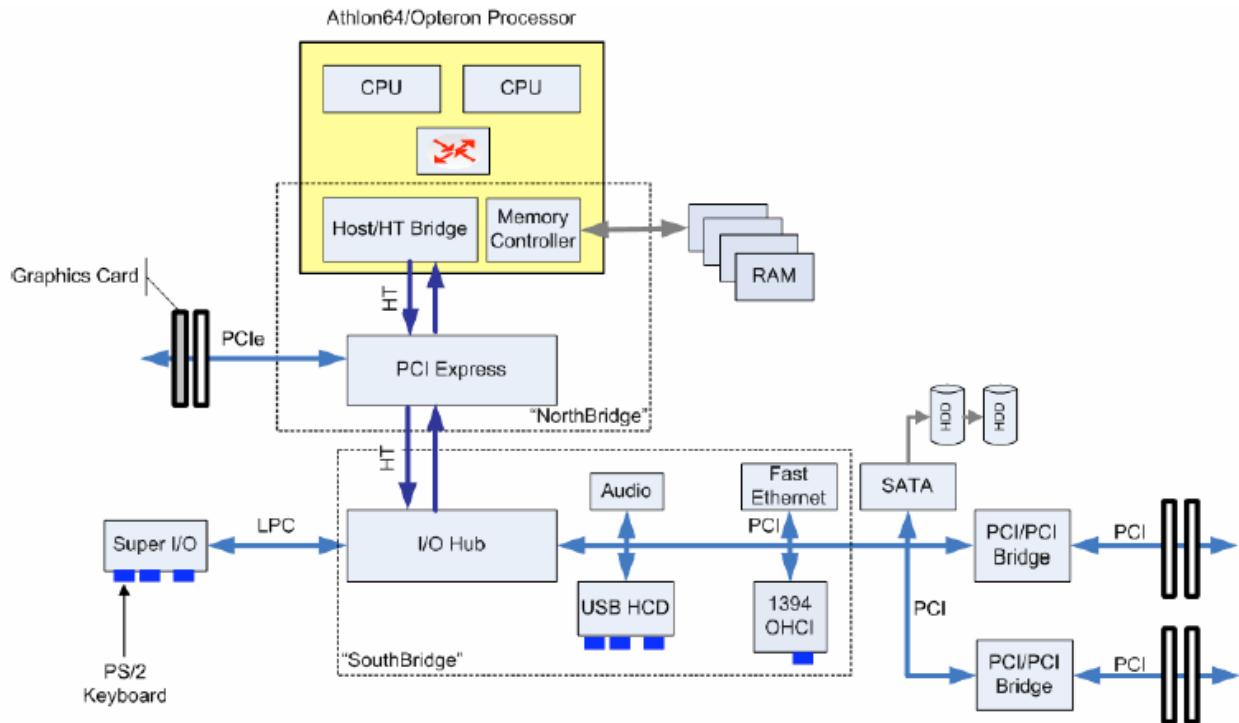


Figure 8: © COSEINC Advanced Malware Labs [x]Device and bus layout for an AMD computer. [12]

Since Firewire can access memory, it should also be possible for it to access other devices through the DMA controller. This would require a transfer through memory and could possibly be detected by the CPU since memory would be altered. Joanna Rutkowska demonstrated how the CPU can remap the memory map in a DMA controller; specifically one of the cases redirected the firewire's (PCI device) requests back towards the PCI bus, causing the system to halt. Given a PCI device can alter these mappings to read memory, it should be possible to find the IO device mapping and remap the DMA controller back at a hard drive's ports. Since the firewire and Serial ATA drives both sit on the same PCI bus, the firewire device might be able to directly access the drive, bypassing the need for using DMA. Current video cards are capable of transferring data directly to each other on a PCI(Express) bus, bypassing the need to communicate back to the CPU, is it a big step to the directly address a hard drive? Could this be another possibly powerful tool in the hands of the forensics investigator?

## Conclusion

Firewire collection presents a few problems:

- Limited availability of firewire ports on computer systems.

- Plugging in a firewire device might require the operating system to activate the port, a slight alteration (artifact) to the state of the system.
- It is possible to crash a system if not done properly; at least the hard drive state would be preserved.
- Primarily, the concepts and tools availability for firewire memory imaging are still immature.

Active memory and live data collection are new to the field of digital forensics and still present a multitude of issues. The desire to collect a snapshot of what is happening in a system is of great value but this must not override the greater value of preserving the integrity of the data collected.



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